

Cellular automata (CA)

- it is a relatively simple tool for modelling complex systems
 - CA often help to understand their behaviour, especially the self-organization property when global order appears spontaneously in systems which are completely defined only with local interactions (examples are flocks of birds, self-replicating life, crystalization etc.)
 - but CA can be used also for modelling in fields of physics where other approaches (such as PDEs) are successfully used, like fluid dynamics, statistical physics etc., especially in systems with more phases or complicated boundary conditions which are changing all the time (deposition etc.)
 - the main advantage of CA is that using simple local rules for interaction of cells (or particles) we can describe many different types of systems (open or isolated, conservative or dissipative etc.)
- note: open and dissipative systems are necessary for self-organization

History of cellular automata (very brief)

- CA were reinvented several times in various fields under different names, but generally as the beginning of CA, it is considered the work of Stanislaw Ulam and John von Neumann in 1940s at Los Alamos NL who proposed a fully discrete dynamical system of cells having internal states described by a finite number of information bits and interacting with nearest neighbours only to simulate self-replicating systems

- CA get wide attention in 1970s when John Conway proposed game of life CA, even with very simple rules it has unexpectedly rich behaviour and it also has the property of universal computation
- in 1980s - Stephen Wolfram studied a family of simple 1D CA rules (now called Wolfram rules) and he noticed that CA is a discrete dynamical system which exhibits many of the behaviours encountered in continuous systems \Rightarrow important for modelling
- later it was also realised that a lattice gas model developed by Hardy, Pomeau and de Pazzis in 1970s (now sometimes called HPP rule) is actually CA, it is used as a ~~as a~~ very simple model in statistical physics \Rightarrow important step in development of the theory of CA
- in 1980s, the first specialized hardware was developed for efficient parallel computations with CA

Basic characteristics of CA

- describes systems discrete in space (lattices) and time
- all states of the system are described by Boolean variables \Rightarrow exact dynamics, no round-off errors
- but only finite set of values for any physical quantity
- simple rules for evolution describing local interactions with neighbouring cells but often nontrivial dynamics

Types of CA

- deterministic or stochastic (probabilistic)
(evolution depends not only on the states of cells at the time t , but also on some random variable(s))
- synchronous (all cells updated at once)
or asynchronous (cells updated sequentially or randomly)
- binary (1 Boolean variable) or more states (more Boolean var.)
- reversible or irreversible (usually)

Definition of deterministic synchronous cellular automaton

- in general, a CA is given by

1) a regular lattice of cells (sites) in a d -dimensional space

2) a set $\Phi(\vec{r}, t) = \{\phi_1(\vec{r}, t), \dots, \phi_m(\vec{r}, t)\}$

of Boolean variables attached to each site \vec{r} of the lattice giving the local state of each cell at times $t=0, 1, 2, \dots$

3) a set of rules $R = \{R_1, R_2, \dots, R_m\}$ specifying the time evolution of the states $\Phi(\vec{r}, t)$ as

$$\phi_j(\vec{r}, t+1) = R_j(\Phi(\vec{r}, t), \Phi(\vec{r} + \vec{\delta}_1, t), \dots, \Phi(\vec{r} + \vec{\delta}_n, t))$$

where $\vec{r} + \vec{\delta}_k$, $k=1, \dots, n$ denote the cells from a given neighbourhood of the cell at \vec{r} .

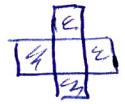
Notes: a) the rule R is identical for all sites (homogeneous) independent of \vec{r} (it depends explicitly only on Φ) and of time t ,
however, spatial and temporal inhomogeneities can be introduced by some $\phi_j(\vec{r}, t)$ systematically equal to 1 in some locations (e.g. boundaries)

to mark cells for which different rules apply,
 or $\phi_j(\vec{r}, t)$ can change with time resulting
 e.g. in different rules for odd or even times

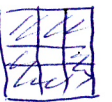
b) this general definition includes also dynamics
 where a new state $\Phi(\vec{r}, t+1)$ is given by
 the states at times $t, t-1, \dots, t-k$
 for we can introduce variables ϕ_j which keep
 necessary information from previous times

c) although the neighbourhood of a cell can be
 in principle arbitrary (but the same for all cells),
 for simplicity and efficiency one usually uses

CAs with von Neumann neighbourhood
 (the nearest cells)

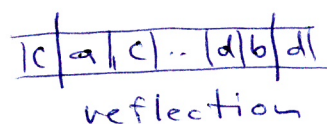
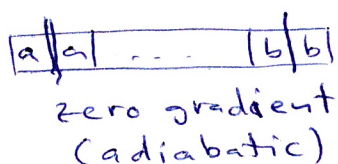
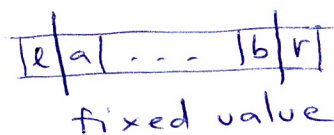
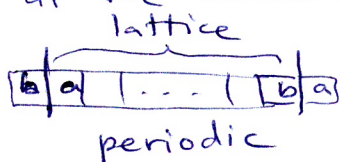


or with Moore neighbourhood
 (the nearest and the second nearest cells)



Boundary conditions

- in simulations, we have to deal only with finite lattices and thus with boundaries
- using additional variables which carry information whether a cell belongs to a specific boundary, we can have different rules at boundaries
- usually, various types of extensions are used at the boundary:



Number of cellular automata

- let us consider a general CA with

q possible states in each cell and

K cells in one neighbourhood (including the cell for which the neighbourhood is defined)

(e.g. for von Neumann neighb. $K = 2d + 1$

and for Moore neighb. $K = 3^d$

where d is the dimension of lattice space),

then there is

$$N_{CA} = q^{q^K} \text{ of different CA}$$

because q^K is the number of states of the whole neighbourhood

and we can assign to each such a state

q possible "new" states of the cell defining the neighbourhood

- for a binary CA: $q = 2$, thus in 1D with $K = 3$ we

$$\text{set } N_{CA} = 2^{2^3} = 256 \text{ cellular automata}$$

(their classification was obtained by

S. Wolfram in *Reviews of Modern Physics* 55, 601 (1983)

see ~~do~~ rules 1-256 in *Mathematica's Cellular Automaton function*)

but in 2D we get

$$\text{von Neumann: } K = 5 \Rightarrow N_{CA} = 2^{2^5} = 2^{32} \approx 4 \cdot 10^9$$

$$\text{Moore: } K = 9 \Rightarrow N_{CA} = 2^{2^9} = 2^{512} \approx 10^{154}$$

we see that the number of CAs grows very fast

thus they are sometimes classified only according

to their typical behaviour if the initial state

is chosen randomly

\Rightarrow 4 basic classes - see examples

in 1D.cellular.Automata.nb

Examples of cellular automata

1) one-dimensional simple CA to model traffic

- it is given by Wolfram's rule 184 (see 10.Cellular.Automata.nb)
- cells ~~have~~ have two possible states
 - 1 - there's a car at that site
 - 0 - there's no car, empty cell
- rules simulate motion of cars in one direction with velocity 1 or 0:
 - ~~the~~ a car can move to the neighbouring right cell if this cell is empty, otherwise it stays
- => "traffic jams" propagate in the opposite direction to cars motion

2) John Conway's Game of Life

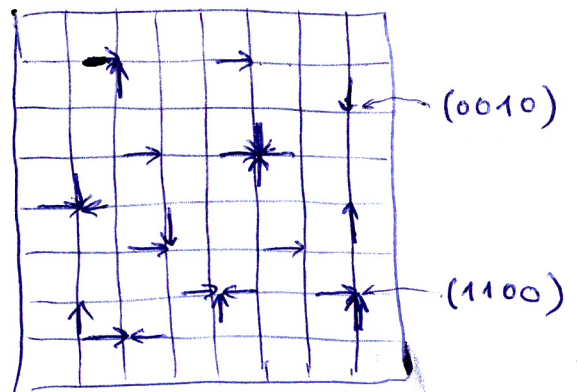
- a well-known CA with simple rules but complex behaviour
- cells have again two possible states - alive (1) or dead (0)
- 2 rules:
 - 1) the cell stays alive if it is surrounded by 2 or 3 live cells, otherwise dies
 - 2) a dead cell becomes alive, if there are exactly 3 live cells in its neighbourhood
- we consider the Moore neighbourhood here (8 surrounding cells) and periodic boundary conditions
- examples can be found at wikipedia or in Game.of.Life.nb

3) Lattice gas model and the HPP rule

- developed in 1970s by Hardy, Pomeau and de Pazzis
but only in 1980s it was recognized as a cellular automaton
- fully discrete dynamics of particles on a lattice (2D, square)
with conservation of momentum and particle number
and in its basic form it is reversible

- the number of particles at one site is limited to
one particle incoming from one direction

(this exclusion principle is actually forced by definition of CA because we cannot describe an arbitrary number of particles in one site within Boolean variables)

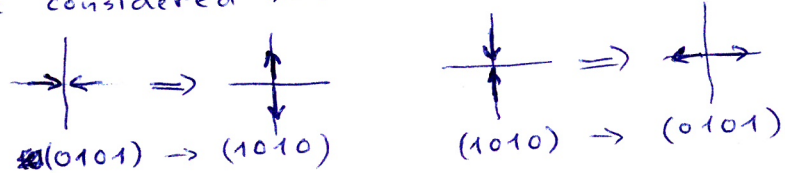


- in the HPP rule, four bits are used to describe incoming particles to a site \vec{r} at the time

e.g. state $s(\vec{r}, t) = (1, 0, 1, 1)$ means there are 3 incoming particles from directions 1, 3, 4
 $\uparrow \rightarrow \downarrow \leftarrow$

- the HPP rule is usually split into two steps

- 1) collisions: determine new directions of particles
if there is a collision, only two collisions are considered (in all other configurations use only step 2)



- 2) motion: after that move all particles along the lattice in their respective direction

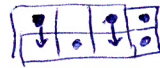
notice that the exclusion principle will be satisfied if the initial conditions were ^{also} chosen to satisfy it and that locally the momentum and particle number are conserved during a collision

- examples of evolution are in lattice Gas, nb where as a boundary condition, the reflection of a particle from the wall is used

4) Sand pile rule

- a simple CA simulating a sand pile or sand motion due to gravity
- again two possible states (except for boundaries)
 - 1 - there's a sand grain in a cell
 - 0 - no grain
- rules in the 2D square lattice:

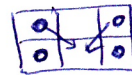
if there is an empty cell below a cell with a grain the grain will fall:



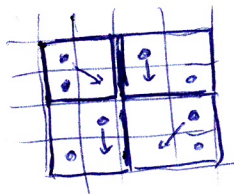
- but a grain can also fall aside



- to avoid conflicts such as

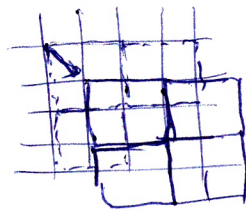


we can use the Margolus neighbourhoods which are changing over the time:



← we solve motion of grains in each 2x2 block independently

and then we shift blocks by 1 to the ~~left~~ ^{right} and down



if we consider boundary condition as a wall in all cells around the lattice

there is no trouble in shifting neighbourhoods

- see Sand.Pile.nb for an example of implementation